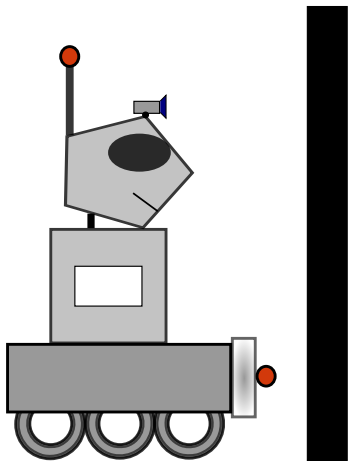
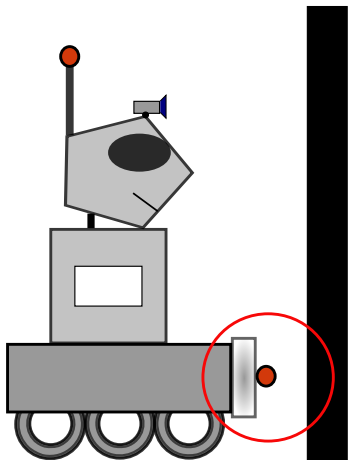


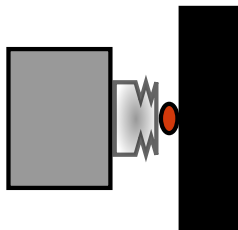
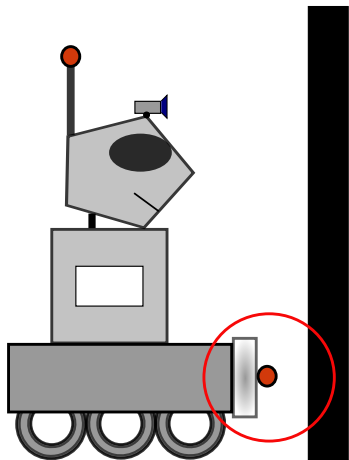
Interactive Timing Analysis for Designing Reactive Systems

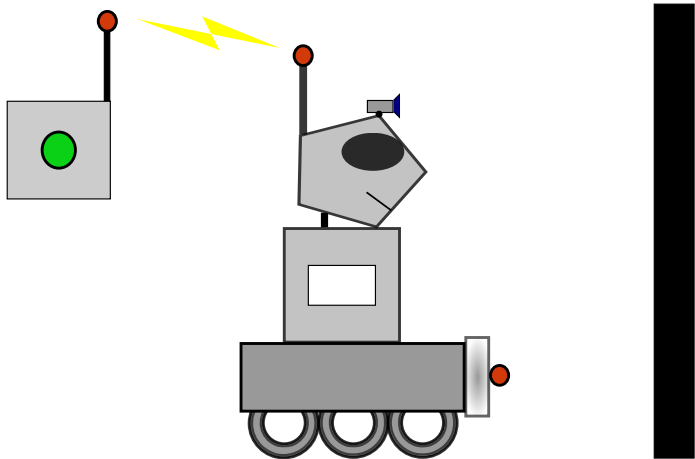
Insa Fuhrmann, David Broman, Reinhard von Hanxleden

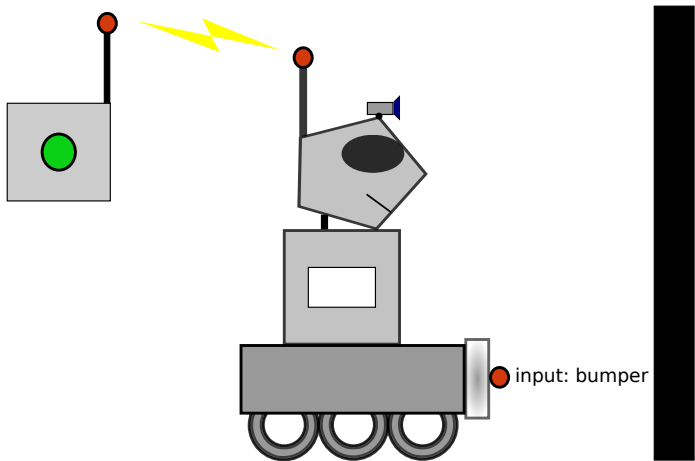
SYNCHRON 2015

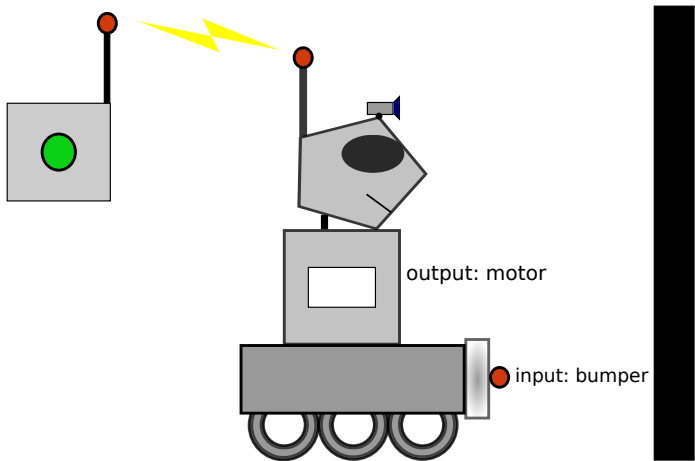


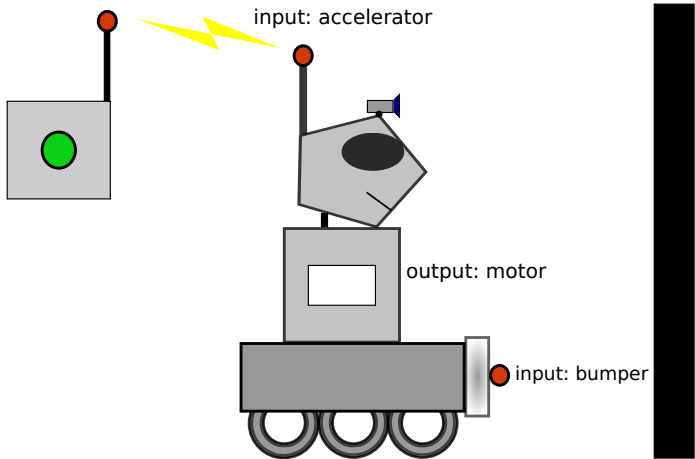


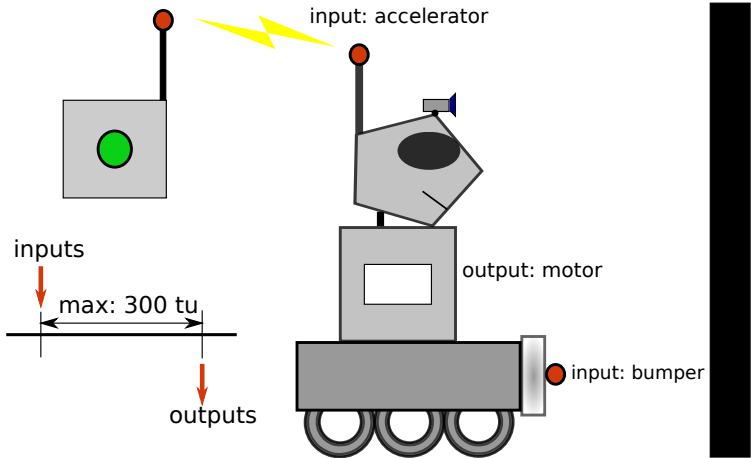


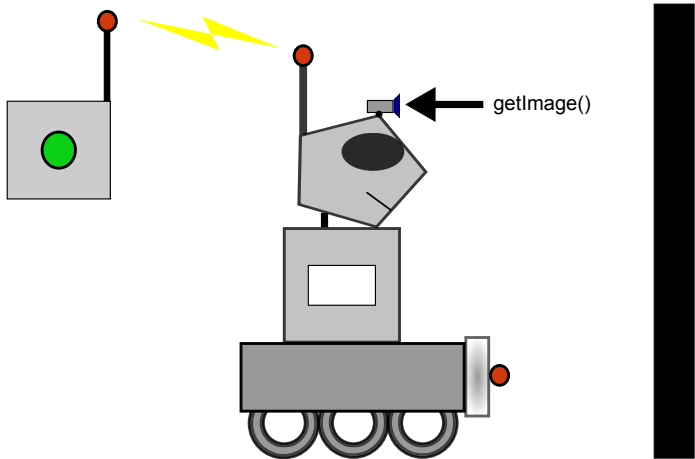


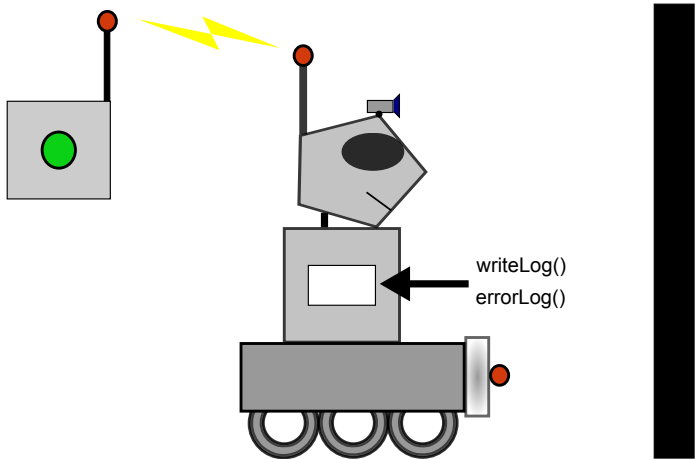


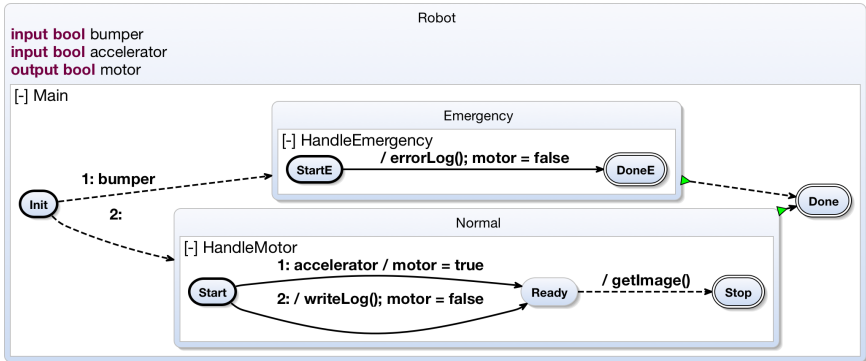








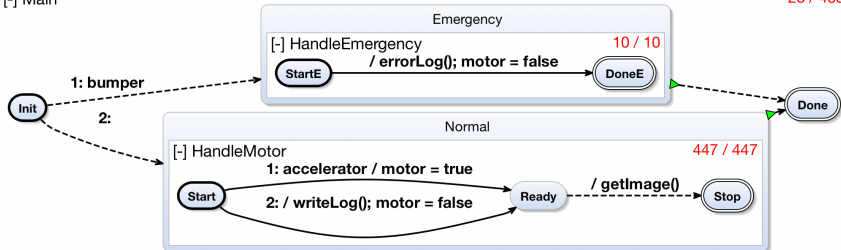




input bool bumper
input bool accelerator
output bool motor

[-] Main

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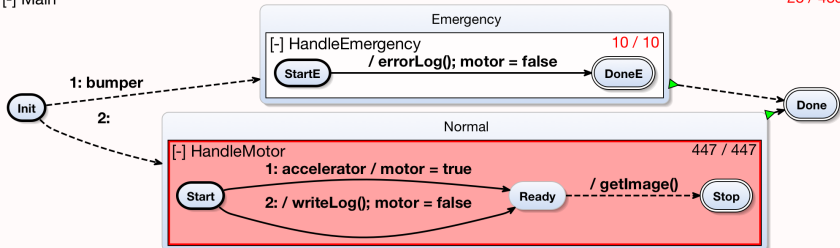
”...hotspots accounted for only 1.2% of the lines of code but contributed 29% of the overall execution time”

G.Bernat et al., Identifying opportunities for worst-case execution time reduction in an avionics system, Ada User Journal, 2007

input bool bumper
 input bool accelerator
 output bool motor

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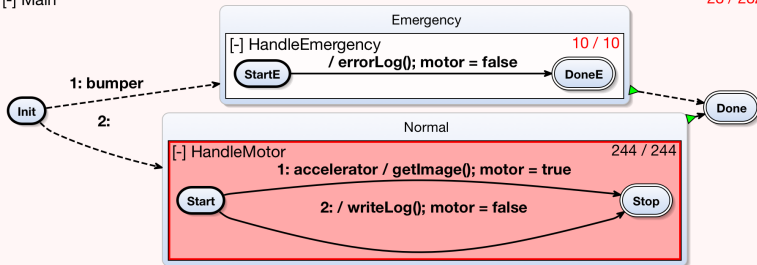
[-] Main

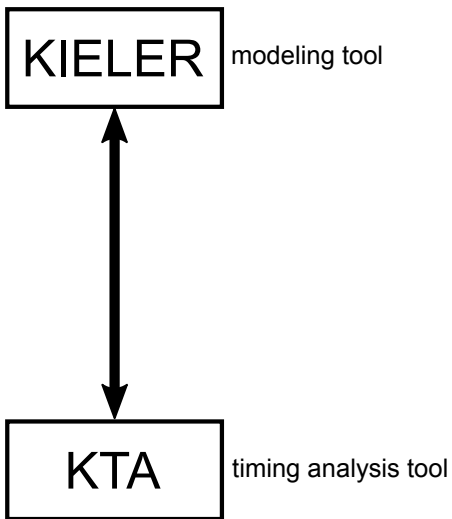


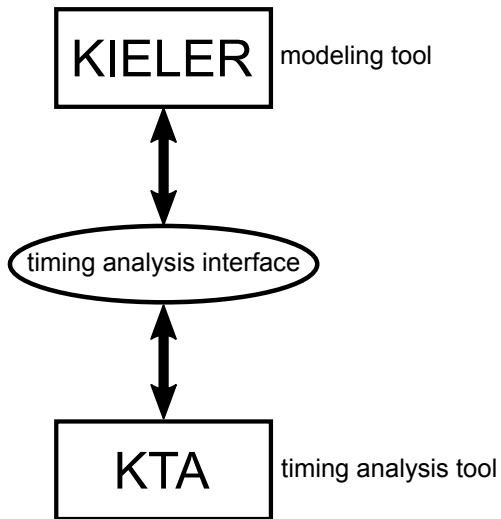
input bool bumper
input bool accelerator
output bool motor
output bool redlight

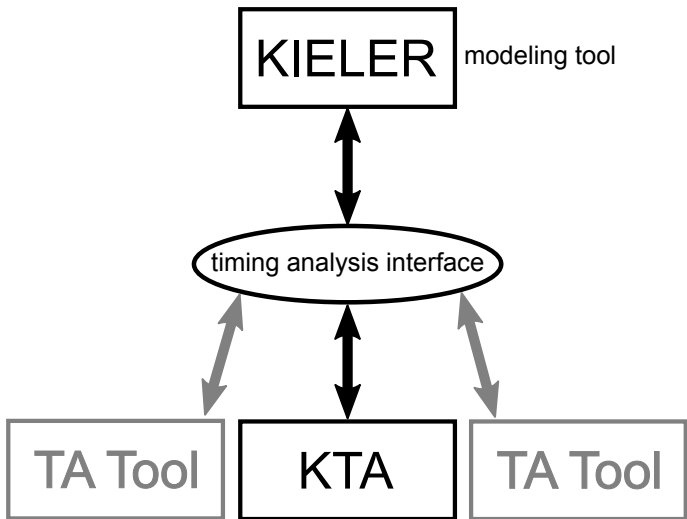
[-] Main

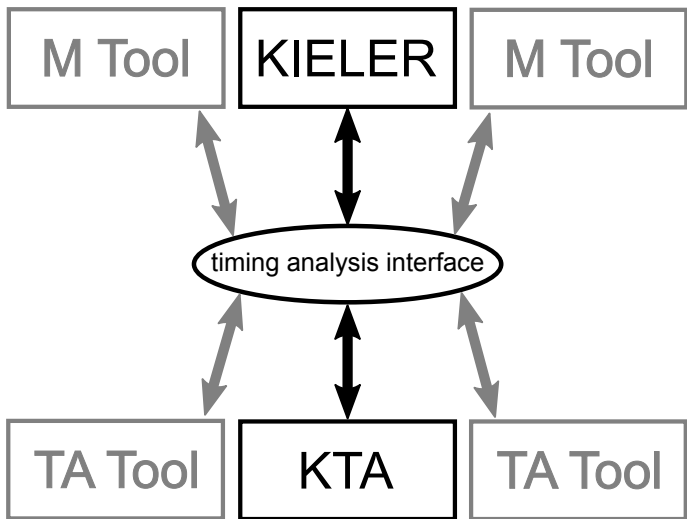
28 / 282

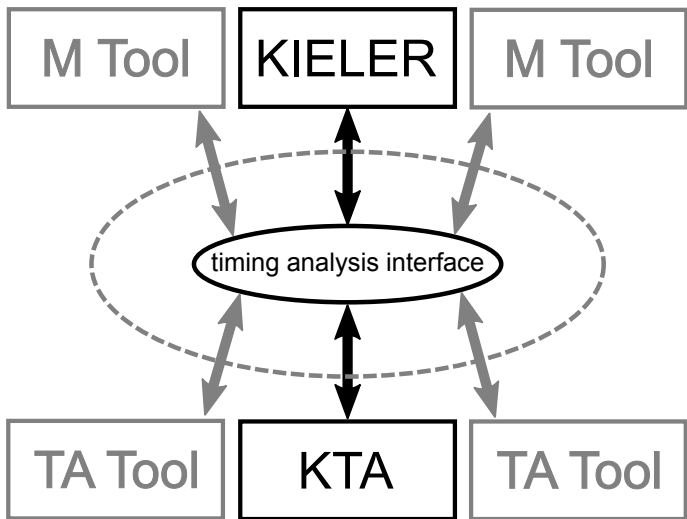


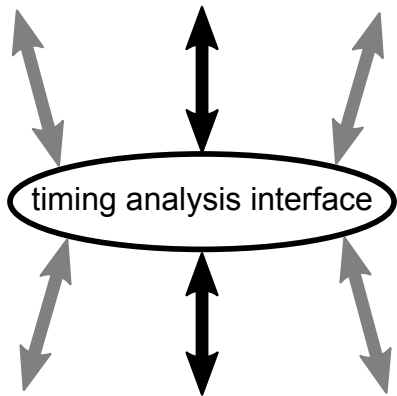




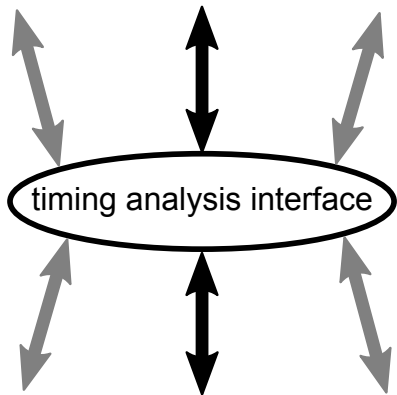




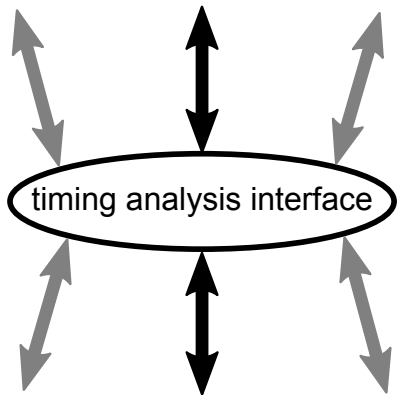




$$t_{req} = ()$$

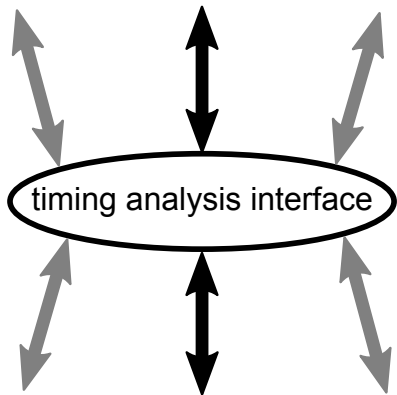


$$t_{req} = (f)$$



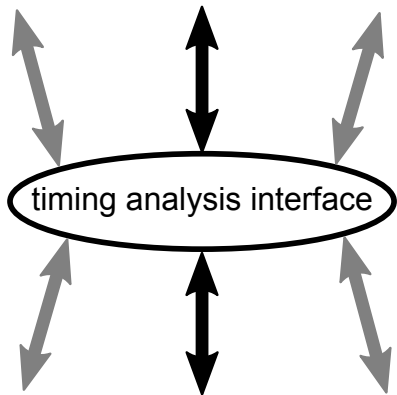
- ▶ Function to be analysed

$$t_{req} = (f, a)$$



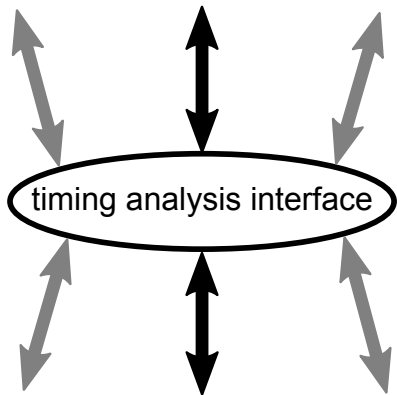
- ▶ Function to be analysed
- ▶ Assumptions on arguments

$$t_{req} = (f, a, g)$$



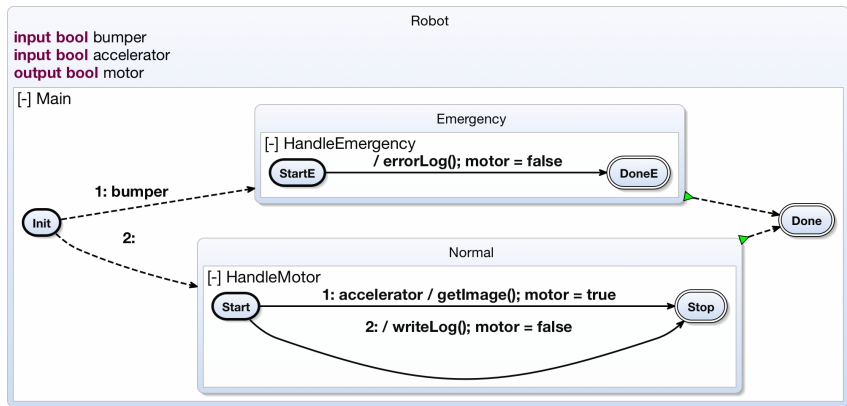
- ▶ Function to be analysed
- ▶ Assumptions on arguments
- ▶ Assumptions on global variables

$$t_{req} = (f, a, g, S)$$



- ▶ Function to be analysed
- ▶ Assumptions on arguments
- ▶ Assumptions on global variables
- ▶ State variables

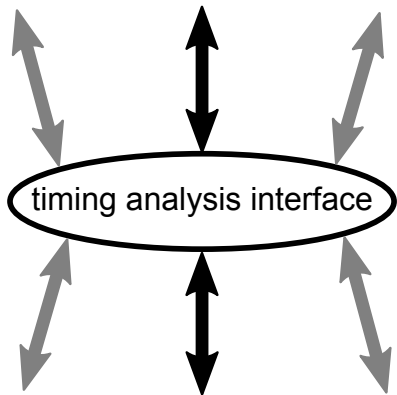
State variables - reactive infeasible paths



State variables - reactive infeasible paths II

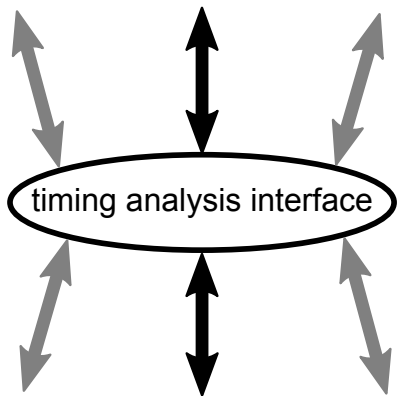
```
1 void tick(){
2     g0 = _GO;
3     g1 = (_GO && bumper);
4     g2 = (PRE_g1);
5     if (g2){
6         errorLog ();
7         motor = 0;}
8     g5 = (PRE_g4);
9     g6 = (g5 && accelerator);
10    if (g6){
11        getImage();
12        motor = 1;}
19    g7 = (g5 && !(accelerator));
20    if (g7){
21        writeLog ();
22        motor = 0;}
23    g3 = (g2 || (g6 || g7));
24    g4 = (_GO && !(bumper));
25    PRE_g1 = g1;
26    PRE_g4 = g4;
27    _GO = 0;
28    return;}
```

$$t_{req} = (f, a, g, S)$$



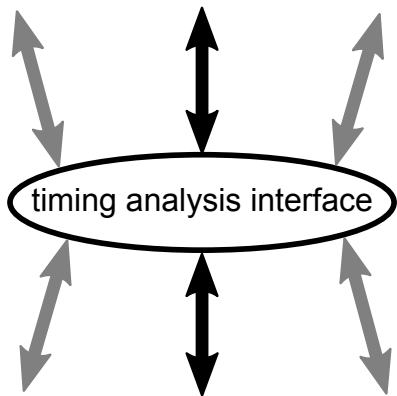
- ▶ Function to be analysed
- ▶ Assumptions on arguments
- ▶ Assumptions on global variables
- ▶ State variables

$$t_{req} = (f, a, g, S, e)$$



- ▶ Function to be analysed
- ▶ Assumptions on arguments
- ▶ Assumptions on global variables
- ▶ State variables
- ▶ Assumptions on called functions

$$t_{req} = (f, a, g, S, e, P)$$



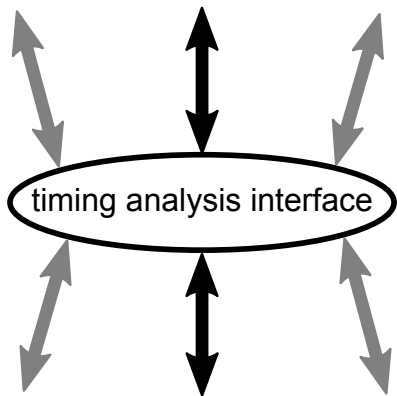
- ▶ Function to be analysed
- ▶ Assumptions on arguments
- ▶ Assumptions on global variables
- ▶ State variables
- ▶ Assumptions on called functions
- ▶ Timing Program Points


```

1 void tick(){
2 //Main
3 //implicit TPP
4     g0 = _GO;
5 //HandleEmergency
6     TPP(1);
7     g1 =(_GO && bumper);
8     g2 =(PRE_g1);
9     if(g2){
10        errorLog ();
11        motor = 0;}
12 //HandleMotor
13     TPP(2);
14     g5 =(PRE_g4);
15     g6 =(g5 && accelerator);
16     if(g6){
17        getImage();
18        motor = 1;}
19        g7 =(g5 && (!(accelerator)));
20        if(g7){
21            writeLog ();
22            motor = 0;}
23 //Main
24     TPP(3);
25     g3 =(g2 || (g6 || g7));
26 //HandleMotor
27     TPP(4);
28     g4 =(_GO && (!(bumper)));
29 //Main
30     TPP(5);
31     PRE_g1 = g1;
32     PRE_g4 = g4;
33     _GO = 0;
34 return ;}
35 //implicit TPP}

```

$$t_{req} = (f, a, g, S, e, P, R)$$



- ▶ Function to be analysed
- ▶ Assumptions on arguments
- ▶ Assumptions on global variables
- ▶ State variables
- ▶ Assumptions on called functions
- ▶ Timing Program Points
- ▶ Analysis Requests

```

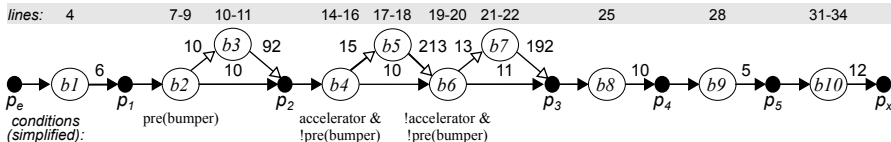
1 void tick(){
2 //Main
3 //implicit TPP
4 g0 = _GO;
5 //HandleEmergency
6 TPP(1);
7 g1 =(_GO && bumper);
8 g2 =(PRE_g1);
9 if(g2){
10 errorLog ();
11 motor = 0;}
12 //HandleMotor
13 TPP(2);
14 g5 =(PRE_g4);
15 g6 =(g5 && accelerator);
16 if(g6){
17 getImage();
18 motor = 1;}

```

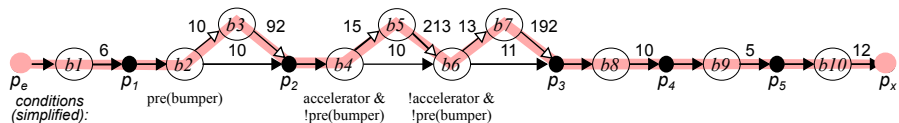
```

19 g7 =(g5 && (!(accelerator)));
20 if(g7){
21 writeLog();
22 motor = 0;}
23 //Main
24 TPP(3);
25 g3 =(g2 || (g6 || g7));
26 //HandleMotor
27 TPP(4);
28 g4 =(_GO && (!(bumper)));
29 //Main
30 TPP(5);
31 PRE_g1 = g1;
32 PRE_g4 = g4;
33 _GO = 0;
34 return;}
35 //implicit TPP}

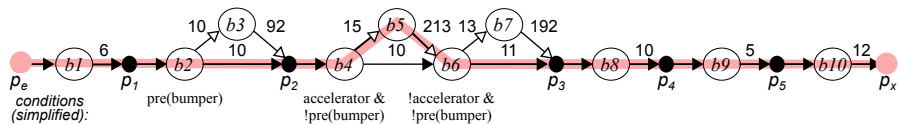
```



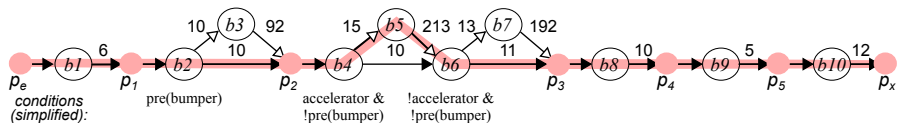
Worst Case Path



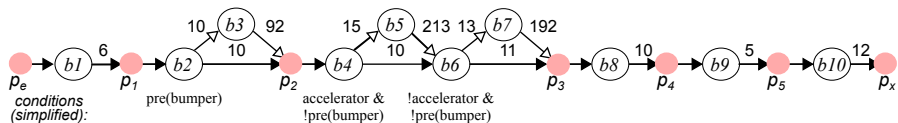
Worst Case Path



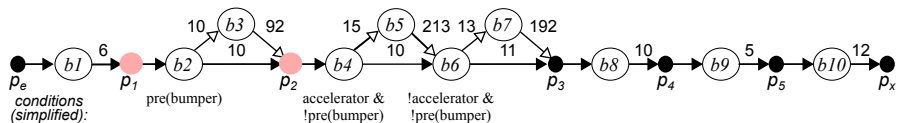
Worst Case Path



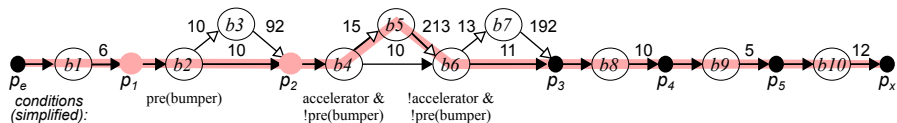
Worst Case Path



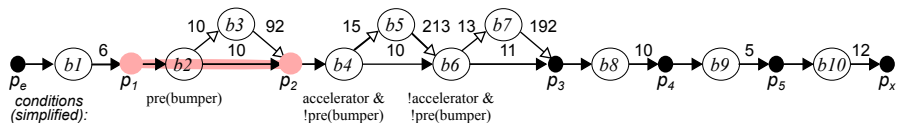
Fractional WCET



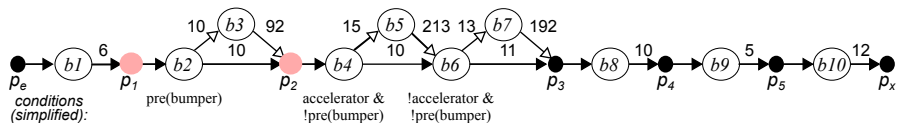
Fractional WCET



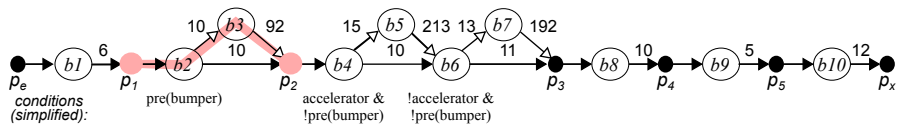
Fractional WCET



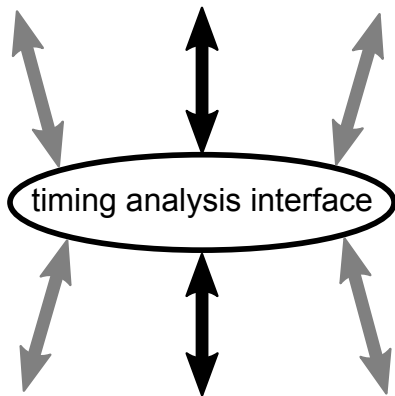
Local WCET



Local WCET

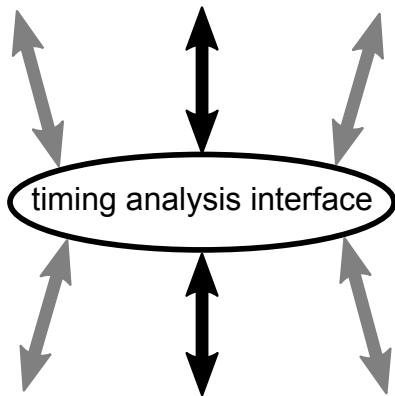


$$t_{req} = (f, a, g, S, e, P, R)$$



- ▶ Function to be analysed
- ▶ Assumptions on arguments
- ▶ Assumptions on global variables
- ▶ State variables
- ▶ Assumptions on called functions
- ▶ Timing Program Points
- ▶ Analysis Requests

$$t_{req} = (f, a, g, S, e, P, R)$$



- ▶ Function to be analysed
- ▶ Assumptions on arguments
- ▶ Assumptions on global variables
- ▶ State variables
- ▶ Assumptions on called functions
- ▶ Timing Program Points
- ▶ Analysis Requests

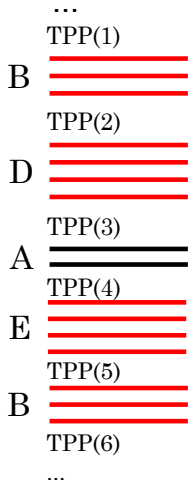
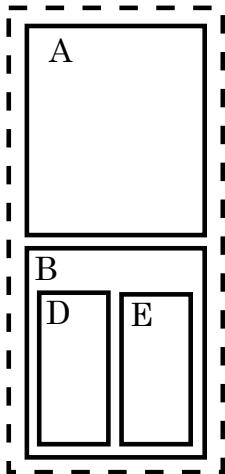
$$t_{res} : R \rightarrow \mathbb{N}_{\perp} \cup \mathcal{P}(\bar{p})$$

Related Work

- Fast WCET Analysis:** Harmon et. al.: Fast, interactive worst-case execution time analysis with back-annotation. Industrial IEEE Transactions on Informatics 2012
- Interactive C-Code analysis:** Ko et. al.: Supporting the specification and analysis of timing constraints. IEEE Real-Time Technology and Applications Symposium 1996
- Analysis of Java Code:** Persson, Hedin: Interactive execution time predictions using reference attributed grammars. WAGA99: 1999
- Matlab/Simulink analysis:** Kirner et. al.: Fully automatic worst-case execution time analysis for Matlab/Simulink models. In: Proceedings of the 14th Euromicro Conference on Real-Time Systems 2002
- SCADE, aiT** Ferdinand et.al.: Combining a high-level design tool for safety-critical systems with a tool for wcet analysis on executables. In: Embedded Real Time Software (ERTS), 2008.

Contributions

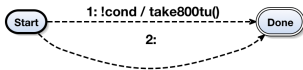
- ▶ Design flow for interactive timing analysis with *hotspot highlighting*, and the separation of *deep*, *flat*, *local* and *fractional* timing values
- ▶ *Timing analysis interface* with separation of concerns of timing analysis for external function calls and for the tick function
- ▶ Concept of Timing Program Points
- ▶ Concept of *reactive infeasible paths*



hierarchicalLWCET

input bool cond
output bool x
output bool y
output bool z
output bool h
output bool complete

[-] A



[-] B

